

Coherence by decreasingness for monoids

Séminaire de réécriture

24 Novembre 2022

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II. Coherence by decreasingness for ARS

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IV. Coherence by decreasingness for \mathcal{K} -SRS

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I. Motivation

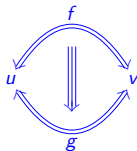
Computing homotopy bases for finitely presented monoids

Let M be a monoid

- ▶ generated by a finite set X of generators,
- ▶ submitted to a finite set $R \subseteq X^* \times X^*$ of oriented relations, called **rewriting rules**.

Question: What is an homotopy basis?

- ▶ A **syzygy**: a relation between relations



- ▶ An **homotopy basis**: a family of syzygies that generates all syzygies of (X, R) .

Convergent presentations of monoids

Let M be a monoid finitely presented.

Question: How to compute an homotopy basis?

Convergent presentation (X, R) of M :

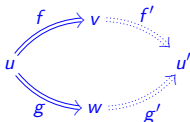
- ▶ **Terminating:** all computations end.
 - ▶ In general define: left degree-wise lexicographic order. Fix an order \prec on elements of X .

$u >_{\text{lex}} v$ iff $\ell(u) > \ell(v)$ or $\ell(u) = \ell(v)$

and $u = x_1 x_2 \dots x_{k-1} x_k x_{k+1} \dots x_n$,

$v = x_1 x_2 \dots x_{k-1} y_k y_{k+1} \dots y_n$ with $y_k \prec x_k$.

- ▶ **Confluent:** the computation converges to the same result (if it exists).



Squier's completion theorem, '94

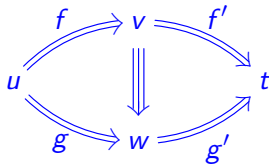
Let $M = \langle X, R \rangle$ be a monoid presented by a finite convergent presentation.

- ▶ A **critical branching**: an overlapping of two rewriting rules that is minimal with respect to the rewriting context.

There are two shapes of critical branchings:



- ▶ The family of syzygies formed by **generating confluences**



indexed by critical branchings, form an homotopy basis.

Another method to compute homotopy bases

Let $M = \langle X, R \rangle$ be a monoid finitely presented such that

- ▶ each word M is represented by a **recurrent form**, i.e. a normal form modulo cycles.
- ▶ each critical pair is **decreasing**, which generalizes the confluence property by adding a well-founded labelling on rewriting steps, [van Oostrom, '94].

Question: How to compute an homotopy basis of (X, R) ?

Example 1: Braid monoid B_3^+

Presentation of B_3^+

$$\langle s, t \mid \alpha : sts \Rightarrow tst \rangle.$$

$$s = \begin{array}{c} \diagdown \quad \diagup \\ \diagup \quad \diagdown \end{array} \quad | \quad t = \begin{array}{c} | \quad \diagdown \\ \diagup \quad | \end{array}$$

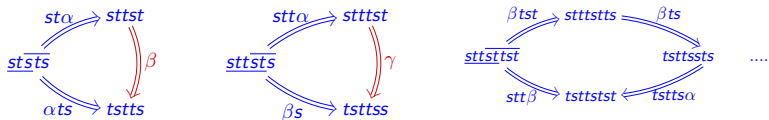
$$\begin{array}{c} \diagdown \quad \diagup \\ \diagup \quad \diagdown \\ \diagdown \quad \diagup \\ \diagup \quad \diagdown \end{array} = \begin{array}{c} \diagup \quad \diagdown \\ \diagdown \quad \diagup \\ \diagup \quad \diagdown \\ \diagdown \quad \diagup \end{array}$$

- ▶ Termination: degree lexicographic order on $s > t$.
- ▶ One non confluent critical branching:

$$\begin{array}{ccc} & \xrightarrow{st\alpha} & sttst \\ \underline{ststs} & & \\ & \xrightarrow{\alpha ts} & tstts \end{array}$$

Example 1: Braid monoid B_3^+

Knuth-Bendix completion: it gives, by adding



a convergent presentation of B_3^+ on the two generators s and t , which is infinite



Theorem. B_3^+ does not admit a finite convergent presentation with the two generators s and t , Kapur & Narendran, '85.

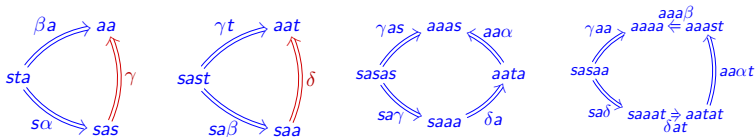
Example 1: Braid monoid B_3^+

Presentation of B_3^+ by adjunction of a new generator a

$$\langle s, t, a \mid \alpha : ta \Rightarrow as, \beta : st \Rightarrow a \rangle$$

standing for the product st .

- ▶ Termination: degree lexicographic order on $s > t > a$.
- ▶ Knuth-Bendix completion:



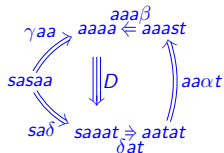
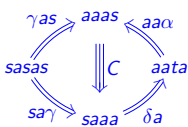
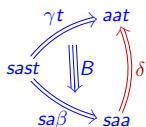
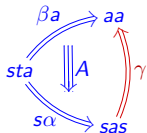
The SRS $\langle s, t, a \mid ta \xrightarrow{\alpha} as, st \xrightarrow{\beta} a, sas \xrightarrow{\gamma} aa, saa \xrightarrow{\delta} aat \rangle$ is a convergent presentation of B_3^+ .

Example 1: Braid monoid B_3^+

By **Squier's completion** of

$$\langle s, t, a \mid \alpha : ta \Rightarrow as, \beta : st \Rightarrow a, \overset{\gamma}{\Rightarrow} aa, saa \overset{\delta}{\Rightarrow} aat \rangle,$$

we obtain an homotopy basis formed by these generating confluences



Another method to compute an homotopy basis for B_3^+

We consider a presentation of B_3^+ , [Alleaume-Malbos, '17]:

$$\Sigma(B_3^+) = \langle s, t \mid \alpha : sts \Rightarrow tst, \beta : tst \Rightarrow sts \rangle.$$

by working on the orientation of the rewriting rules instead of adding new generators.

Question:

Can we compute an homotopy basis of $\Sigma(B_3^+)$?

Example 2: Plactic monoid P_n

Plactic monoid P_n

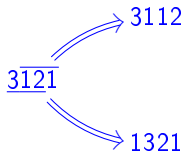
- ▶ generated by $1 < 2 < \dots < n$
- ▶ and submitted to Knuths relations

$$x \leq y < z, \quad yzx \Rightarrow yxz$$

$$x < y \leq z, \quad zxy \Rightarrow xzy$$

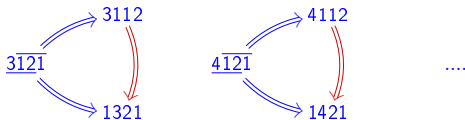
with $x, y, z \in \{1, 2, 3, \dots, n\}$.

- ▶ Termination: degree lexicographic order on $1 < 2 < \dots < n$.
- ▶ Some non confluent critical branchings:



Example 2: Plactic monoid P_n

Knuth-Bendix completion: it gives, by adding



a convergent presentation of P_n on the generators $1, \dots, n$, which is infinite.

Theorem. $\forall n > 3$, P_n does not admit a finite convergent presentation by Knuth-Bendix completion of the Knuth presentation with the degree lexicographic order, *Kubat and Okninski, '14*.

Example 2: plactic monoids P_n

- ▶ **Young tableau** over $\{1, \dots, n\}$:

2	2	3	5	5
3	3	4	6	
4	5			

$Col(n)$: set of columns.

- ▶ **Schensted algorithm:**

any word w in $\{1, \dots, n\}^*$ \rightsquigarrow Young tableau $P(w)$.

- ▶ Notation: any two columns $u \times v$ iff $T = c_u c_v$ is not a Young tableau.
- ▶ Oriented relations on $Col(n)^*$:

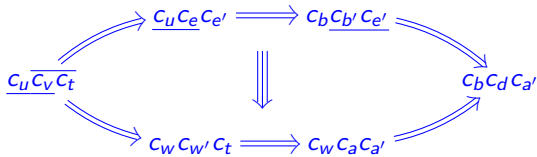
$$Col_2(n) = \{ c_u c_v \Rightarrow c_w c_{w'} \mid u \times v \in Col(n) \text{ and } P(uv) = c_w c_{w'} \}$$

Example 2: Plactic monoid P_n

- ▶ The **column presentation** $Col_2(n)$ is a convergent presentation of P_n .
- ▶ By **Squier's completion** of

$$Col_2(n) = \{ c_u c_v \Rightarrow c_w c_{w'} \mid u^x v \in Col(n) \text{ and } P(uv) = c_w c_{w'} \},$$

we obtain an homotopy basis, denoted $Col_3(n)$, formed by these confluences



indexed by critical branchings, [Hage-Malbos, '16].

Another method to compute an homotopy basis for P_n

We consider a presentation of P_n :

$$\Sigma(P_n) = \langle 1, \dots, n \mid yzx \xrightarrow{\alpha_{x,y,z}} yxz, x \leq y < z; \quad xzy \xrightarrow{\beta_{x,y,z}} zxy, x < y \leq z \rangle.$$

by working on the orientation of the rewriting rules instead of adding new generators.

Questions:

Can we compute an homotopy basis of $\Sigma(P_n)$? Is it smaller than the homotopy basis $Col_3(n)$ of the column presentation $Col_2(n)$?

II. Coherence by decreasingness for ARS

One-dimensional polygraphs

A **1-polygraph** X : modelization of abstract rewriting system.

- ▶ set of **0**-generators X_0 ,
- ▶ set of **1**-generators X_1 , **called rewriting steps**,
- ▶ source and target maps

$$X_0 \begin{array}{c} \xleftarrow{s_0} \\ \xleftarrow{t_0} \end{array} X_1.$$

Free categories over a 1-polygraph X :

- ▶ **free 1-category** X_1^* whose morphisms:

$$x_0 \xrightarrow{u_1} x_1 \xrightarrow{u_2} x_2 \xrightarrow{u_3} \dots \xrightarrow{u_n} x_n$$

with $u_i \in X_1$.

- ▶ **free (1,0)-category** X_1^\top whose morphisms:

$$x_0 \xleftarrow{u_1} x_1 \xleftarrow{u_2} x_2 \xleftarrow{u_3} \dots \xleftarrow{u_n} x_n$$

with $u_i : x_{i-1} \rightarrow x_i$ or $u_i : x_i \rightarrow x_{i-1}$.

van Oostrom's decreasingness theory: well-founded labelling

A **Well-founded labelled 1-polygraph** is a data (X, I, \leq_I, ψ) made of:

- ▶ a 1-polygraph X ,
- ▶ well-founded ordered set (I, \leq_I) of labels,
- ▶ **well-founded labelling**:

$$\begin{aligned} \psi : X_1 &\longmapsto (I, \leq_I) \\ u &\longmapsto \psi(u). \end{aligned}$$

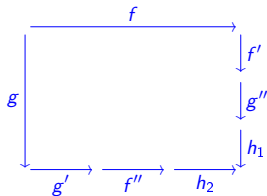
Given a 1-cell $f = x_0 \xrightarrow{u_1} x_1 \xrightarrow{u_2} x_2 \xrightarrow{u_3} \dots \xrightarrow{u_n} x_n \in X_1^*$, we denote by

$$L^I(f) = \{\psi(u_1), \dots, \psi(u_n)\}$$

the set of labels of rewriting steps in f .

Locally decreasingness

- ▶ A **local branching** (f, g) : a pair of rewriting steps f and g with $s_0(f) = s_0(g)$.
- ▶ A **decreasing confluence diagram** of (f, g) is defined by



with $f', g', f'', g'', h_1, h_2 \in X_1^*$ and such that

- ▶ $k < \psi(f)$, for all k in $L^1(f')$,
- ▶ g'' is an identity or a 1-generator labelled by $\psi(g') = \psi(g)$,
- ▶ $k < \psi(f)$ or $k < \psi(g)$, for all k in $L^1(h_1)$.
- ▶ symmetrically for g', f'' and h_2 .

van Oostrom's confluence theorem

Let (X, I, \leq_I, ψ) be a well-founded labelled 1-polygraph.

- ▶ A **(finite) multiset** over I : a function $A : I \rightarrow \mathbb{N}$ such that the set $\{i \in I \mid A(i) \neq 0\}$ is finite.
 - ▶ We denote $M(I)$ the set of finite multisets over I .
 - ▶ **Rq**: we generalize \leq_I to a well-founded order \leq_{mul} on $M(I)$.
- ▶ **lexicographic maximum measure**: the multiset of the lexicographically maximal step labels [van Oostrom, '94].

$$\begin{array}{lcl} |\cdot| : X_1^* & \longmapsto & M(I) \\ f & \longmapsto & |f| \end{array}$$

- ▶ **Theorem**: Any locally decreasing well-founded polygraph is confluent, [van Oostrom, '94].

Recurrent forms, [Zilli, '84]

Let X be a 1-polygraph.

- ▶ The **X -congruence**: equivalence relation on X_0 defined by

$$x \approx^X y \iff \exists f : x \rightarrow y \in X_1^T.$$

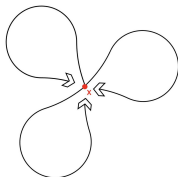
We denote

- ▶ Equivalence set: $\bar{X} = X_0 / \approx^X$.
- ▶ Canonical projection: $\pi : X_0 \rightarrow \bar{X}$.
- ▶ A **recurrent 0-generator** x of a 1-polygraph X : for each 1-cell

$$f : x \rightarrow y \in X_1^*,$$

there exists a 1-cell

$$g : y \rightarrow x \in X_1^*.$$



Recurrence property

- ▶ A **section** of a 1-polygraph X : a section of π

$$\begin{array}{l} r : \bar{X} \longmapsto X_0 \\ \bar{x} \longmapsto r(\bar{x}), \end{array} \quad \pi(r(\bar{x})) = \bar{x}.$$

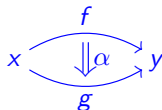
- ▶ We denote $r(\bar{x}) = \tilde{x}$.
 - ▶ a **recurrent section** r : for each $x \in X_0$, \tilde{x} is recurrent.
- ▶ A **recurrent 1-polygraph** X : admits a recurrent section
 - ▶ each equivalence class \bar{x} is represented by a recurrent 0-generator \tilde{x} .

coherent extensions of 1-polygraphs

- ▶ A **sphere** in X_1^\top : a pair (f, g) of 1-cells such that

$$s_0(f) = s_0(g) \quad \text{and} \quad t_0(f) = t_0(g).$$

- ▶ Denote by $\text{Sph}(X_1^\top)$ the set of spheres of X_1^\top .
- ▶ A **cellular extension** of X_1^\top : a set X_2 with a map $X_2 \rightarrow \text{Sph}(X_1^\top)$.
 - ▶ Elements of X_2 are called **2-generators**



- ▶ A **(2, 0)-polygraph** X is defined by
 - ▶ a 1-polygraph (X_0, X_1) ,
 - ▶ a cellular extension X_2 of X_1^\top

coherent extensions of 1-polygraphs

Let X be a $(2, 0)$ -polygraph.

- ▶ **Free $(2, 0)$ -category X_2^\top** on X .
 - ▶ its 2-cells are generated via \star_0 and \star_1 of 2-generators in X_2 , of their inverses and of identities of 1-cells of X_1^\top .
- ▶ An **homotopy basis X_2** : a cellular extension of X_2^\top such that for any sphere (f, g) , there exists a 2-cell

$$A : f \Rightarrow g \in X_2^\top$$

- ▶ A **coherent $(2, 0)$ -polygraph X** is defined by
 - ▶ a 1-polygraph (X_0, X_1) ,
 - ▶ an homotopy basis X_2 .

Coherence properties

A **coherently recurrent (2, 0)-polygraph** is defined by

- ▶ a recurrent 1-polygraph (X_0, X_1) ,
- ▶ a recurrent section $r(\bar{x}) = \tilde{x}$ of (X_0, X_1) ,
- ▶ a **recurrent cellular extension** X_2 of X_1^\top : for any 1-cell

$$f : \tilde{x} \rightarrow \tilde{x} \in X_1^*,$$

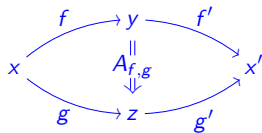
there exists a 2-cell

$$A : f \Rightarrow 1_{\tilde{x}} \in X_2^\top.$$

Coherence properties

A **coherently locally decreasing** $(2, 0)$ -polygraph:

- ▶ a well-founded labelled 1-polygraph (X_0, X_1) ,
- ▶ a well-founded labelling ψ of (X_0, X_1) ,
- ▶ a **decreasing cellular extension** X_2 of X_1^T : for any local branching (f, g) , there exists a decreasing confluence diagram



and a 2-cell $A_{f,g}$ in X_2^T .

Main theorem

Theorem. Let X be a coherently recurrent $(2,0)$ -polygraph.
If X is locally coherently decreasing, then X is coherent.

III. Coherence by decreasingness for SRS

Application: braid monoid B_3^+

- ▶ **Presentation** of B_3^+

$$\Sigma(B_3^+) = \langle s, t \mid \alpha : sts \Rightarrow tst, \beta : tst \Rightarrow sts \rangle.$$

- ▶ **Rule-labelling** of $\Sigma(B_3^+)$, [van Oostrom, '08].

- ▶ well-founded labelling of **the set of rewriting steps** Σ_{step}

$$\psi : \Sigma_{step} \mapsto (\mathbb{N}, \leq),$$

- ▶ invariance with the **algebraic context**

$$\psi(u\alpha v) = \psi(\alpha) \quad \text{and} \quad \psi(u\beta v) = \psi(\beta), \quad u, v \in \{s, t\}^*,$$

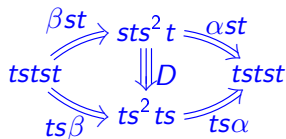
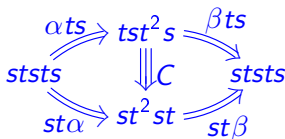
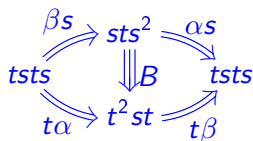
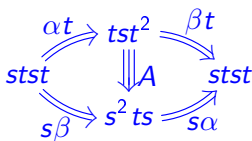
- ▶ **constant rule-labelling**

$$\psi(u\alpha v) = \psi(\alpha) = 1 \quad \text{and} \quad \psi(u\beta v) = \psi(\beta) = 1, \quad u, v \in \{s, t\}^*.$$

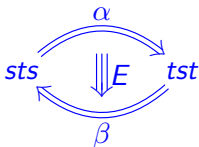
Application: braid monoid B_3^+

Extension of $\Sigma(B_3^+)$ by

- ▶ generating decreasing confluence diagrams



- ▶ a generating cycle



Application: braid monoid B_3^+

- ▶ The **extended presentation** of B_3^+

$$\langle s, t \mid \alpha : sts \Rightarrow tst, \beta : tst \Rightarrow sts \mid A, B, C, D, E \rangle$$

is **coherent**.

- ▶ By **homotopical reduction**:

$$\langle s, t \mid \alpha, \emptyset \mid \cancel{A}, \cancel{B}, \cancel{C}, \cancel{D}, \cancel{E} \rangle.$$

with $A > B > C > D > E$ and $\beta > \alpha$.

- ▶ A coherent presentation of B_3^+

$$\langle s, t \mid \alpha : sts \Rightarrow tst \mid \emptyset \rangle.$$

IV. Coherence by decreasingness for \mathcal{K} -SRS

Application: plactic monoid P_n

\mathcal{K} -Presentation $\Sigma(P_n)$ of P_n

▶ a set $[n] = \{1, 2, \dots, n\}$ of generators,

▶ a **crystal basis** on $[n]$

$$1 \xrightarrow{1} 2 \xrightarrow{2} \dots \xrightarrow{n-2} n-1 \xrightarrow{n-1} n,$$

▶ a set of oriented Knuth's relations

$$x \leq y < z, \quad \alpha_{x,y,z} : yzx \Rightarrow yxz$$

$$x < y \leq z, \quad \beta_{x,y,z} : xzy \Rightarrow zxy$$

with $x, y, z \in [n]$.

Application: compatibility with Kashiwara's operators

Kashiwara's operators on $[n]$

- ▶ For each i in $\{1, 2, \dots, n-1\}$

$$f_i(i) = i + 1 \quad \text{and} \quad e_i(i + 1) = i$$

Kashiwara's operators on $[n]^*$

- ▶ Example: compute $f_1(121)$

$$\begin{array}{ccc} 1 & 2 & 1 \\ + & - & + \\ & & + \end{array}$$

Then, $f_1(121) = 12f_1(1) = 122$.

Highest weight word w in $[n]^*$: $e_i(w)$ not defined for each $i \in \{1, \dots, n-1\}$.

Application: compatibility with Kashiwara's operators

- ▶ **Remark:** the image of a rewriting rule by Kashiwara's operators is a rewriting rule in the same direction.

Example: Let $\alpha_{1,1,2} : 121 \Rightarrow 112$. We have:

- ▶ $f_1(121) = 122$ and $f_1(112) = 212$,
- ▶ $\beta_{1,2,2} : 122 \Rightarrow 212$,

Then,

$$f_1(\alpha_{1,1,2}) = \beta_{1,2,2} : f_1(121) \Rightarrow f_1(112).$$

- ▶ We obtain a **\mathcal{K} -2-polygraph**, [Uran, '22], presenting P_n .
- ▶ **Remark:** The Young tableaux are recurrent forms.

Application: Well-founded labelling of $\Sigma(P_n)$

A **well-founded labelling** on the set of rewriting steps Σ_{step}

$$\psi : \Sigma_{step} \mapsto (\mathbb{N}, \leq),$$

► a **totally ordered set of labels** on the set of rewriting rules Σ_2

$$\psi : \Sigma_2 \mapsto \mathbb{N}$$

► For $wfw' \in \Sigma_{step}$, with $w, w' \in [n]^*$ and $f \in \Sigma_2$

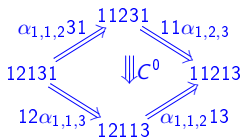
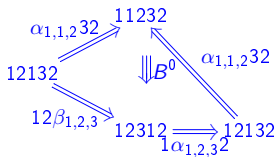
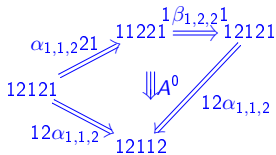
$$\psi(wfw') = (|w| - 1, \psi(f)),$$

with $|1_{[n]^*}| = 0$ and $|w|$ indicates the length of w .

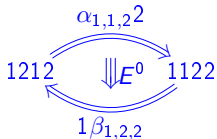
Application: plactic monoid P_n

Extension of $\Sigma(P_n)$ by

- ▶ generating \mathcal{K} -decreasing hw-confluence diagrams



- ▶ a generating hw-cycle



Application: plactic monoid P_n

- ▶ The **extended presentation** of P_n

$$\langle 1, \dots, n \mid \alpha_{1,1,2}, k(\alpha_{1,1,2}) \mid A^0, k(A^0), B^0, k(B^0), C^0, k(C^0), E^0, k(E^0) \rangle.$$

with k a **a sequence of Kashiwara's operators**, is coherent.

- ▶ By **homotopical reduction**:

$$\langle 1, \dots, n \mid \alpha_{1,1,2}, k(\alpha_{1,1,2}) \mid \cancel{A^0}, \cancel{k(A^0)}, \cancel{B^0}, \cancel{k(B^0)}, C^0, k(C^0), E^0, k(E^0) \rangle.$$

- ▶ A coherent presentation of P_n

$$\langle 1, \dots, n \mid \alpha_{1,1,2}, k(\alpha_{1,1,2}) \mid C^0, k(C^0), E^0, k(E^0) \rangle.$$

with k a **a sequence of Kashiwara's operators**

V. Perspectives

Perspectives

- ▶ Applying our result of coherence by decreasingness to other algebraic structures with quasi-terminating presentations.
- ▶ apply decreasingness techniques to study Chinese monoids, [Endrullis-Klop, '19].